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Reachability of eigenspaces for interval circulant matrices in max-algebra

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Abstract

A nonnegative matrix A is said to be strongly robust if its max-algebraic eigencone is universally reachable, i.e., if the orbit of any initial vector ends up with a max-algebraic eigenvector of A. Consider the case when the initial vector is restricted to an interval and A can be any matrix from a given interval of nonnegative circulant matrices. The main aim of this paper is to classify and characterize the six types of interval robustness in this situation. This naturally leads us also to study the max-algebraic spectral theory of circulant matrices and the relation of inclusion between attraction cones of circulant matrices in max-algebra.

Keywords: Max-algebra, circulant matrices, interval analysis, reachability. *AMS classification: 15A18, 15A80, 65G40, 93C55*

1. Introduction

Max-algebra has applications in such fields as discrete event systems and scheduling theory (among others) [2, 4, 11], and plays a crucial role in the study of discrete event systems in connection with optimization problems such as scheduling or project management in which the objective function depends on the maximum and times operations (or equivalently maximum and plus via a logarithmic transform). Notice that the main principle of

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discrete events systems consisting of n entities is that the entities work interactively, i.e., a given entity must wait before proceeding to its next event until certain others have completed their current events. The steady states of such systems correspond to the max-algebraic eigenvectors of the matrices that describe them, therefore the investigation of reachability of the set of eigenvectors from a given state by a given system is important for such applications. Matrices for which the steady states of the corresponding systems are reached with any nontrivial starting vector are called robust, see [4] Section 8.6.

In practice, matrix entry values are not exact numbers and usually are contained within intervals, and therefore interval arithmetic is an efficient way to represent matrices in a guaranteed way on a computer. A maxalgebraic (tropical) version of interval analysis was developed, e.g., in [12], which emphasized the polynomiality of some algorithms of max-algebraic interval analysis. That polynomiality was in striking contrast with NPhardness of relevant algorithms previously known in usual interval analysis. Independently, [7] developed a theory of some max-algebraic linear systems with interval coefficients and optimization problems over such systems.

When developing interval extensions of linear algebra problems a whole range of solvability problems routinely arises, by considering all possible combinations of quantifiers (as in Definition 2.8 of the present paper). In classical linear algebra this leads to the notions of united solutions, controllable solutions and tolerable solutions [17, 20]. In max-algebra we similarly have, e.g., four types of interval extensions of the max-algebraic spectral problem [8] or two types of interval extensions of robustness studied in [15].

Similarly to [15], the present paper also considers max-algebraic interval extensions of robustness and reachability problems. However, we focus on matrices of a certain special type: circulants. In usual algebra, circulant matrices have a number of geometric applications [6]. A more recent application of circulants can be found in [21]. There, an algebraic construction based on circulant matrices allows for designing LDPC codes with efficient encoder implementation, in contrast to designing LDPC codes based on random construction techniques which make it difficult to store and assess a large parity-check matrix or to analyze the performance of the code. In max-algebra, circulant matrices appear to describe the periodic regime of sequences of matrix powers [4, 18]. It is also easy to see that circulant matrices of a given dimension form a commutative semigroup, both in max-algebra and in usual linear algebra.

When considering matrices of special type, it is natural to require that the set of matrices that is an interval extension of such a matrix can contain matrices of that type only. This is a basic idea behind the notion of interval circulant matrix defined here. The main aim of the present paper is thus to classify and characterize the six types of interval robustness for circulant matrices in max-algebra. However, obtaining such a characterization is not possible without a deeper study of properties of circulant matrices in maxalgebra, which is itself of some theoretical interest.

We now outline the organization of the paper and the results obtained there. Section 2 is devoted to some basic notions of max-algebra and its connections to the theory of digraphs and max-algebraic convexity. In particular, we revisit the max-algebraic spectral theory here, focusing on the eigencone and the attraction cone associated with an arbitrary eigenvalue, the cyclicity of critical graphs and the ultimate periodicity of max-algebraic matrix powers and orbits.

Section 3 presents some known as well as some new results on the spectral theory and attraction cones of circulant matrices. In particular, Proposition 3.7 describes the critical node sets of circulant matrices and presents several formulae for the cyclicity of the critical graph of a circulant matrix. This result combines together some facts that have been previously obtained or stated in [14, 15, 22]. The main new result of this section is Theorem 3.10, which deals with a particular problem of inclusion of the attraction cones of circulant matrices A and B satisfying $A \leq B$ and having the same maximum cycle mean. It appears that inclusion $\operatorname{attr}(A) \subseteq \operatorname{attr}(B)$ holds for such circulant matrices. Note that it does not hold for general matrices, as Example 2.24 demonstrates. Section 3 also contains several motivating examples. The proofs of Proposition 3.7 and Theorem 3.10 are deferred to Section 5.

Based on the result about inclusion of attraction cones of Theorem 3.10, Section 4 characterizes various types of interval robustness which are described in Definition 2.8. Some of them can be verified in polynomial time, see Theorems 4.7, 4.9, 4.15. Other types of robustness reduce to max-algebraic two-sided systems of equations and inequalities for which efficient algorithms exist but the problem of constructing a polynomial algorithm remains open. See Theorems 4.11, 4.13, 4.14.

Subsection 5.1 presents a proof of Proposition 3.7. The proof uses the fact that any circulant matrix is strictly visualized in the sense of [19] and relies in part on the results of [9, 10].

Subsection 5.2 presents a proof of Theorem 3.10. In particular, the proof

draws upon the role of cyclic classes in the max-linear systems of equations describing attraction cones, as presented in [4] Chapter 8 and [18].

2. Preliminaries

2.1. Main definitions and problem statements

By max-algebra we mean the set of nonnegative numbers \mathbb{R}_+ equipped with the usual multiplication $a \cdot b$ and the idempotent addition $a \oplus b :=$ max(a, b). These arithmetical operations are then routinely extended to matrices and vectors: in particular, $(A \otimes B)_{i,k} = \bigoplus_j A_{i,j} \cdot B_{j,k}$ and $(A \oplus B)_{i,j} =$ $A_{i,j} \oplus B_{i,j}$ for any two nonnegative matrices A, B of appropriate sizes. We will also consider the max-algebraic powers of matrices $A^k := \underline{A \otimes \ldots \otimes A}$.

In what follows, we will be interested in the orbits of vectors under the action of matrices, that is, the sets

$$\mathcal{O}(A, x) = \{x, A \otimes x, A^2 \otimes x, \ldots\},\tag{1}$$

and especially in the case when the orbit of a vector hits an eigenvector of A. Let us now give formal definitions related to the max-algebraic eigenproblem.

Definition 2.1 (Eigenvalues and Eigenvectors). A value $\lambda \in \mathbb{R}_+$ is called a *(max-algebraic) eigenvalue* of $A \in \mathbb{R}^{n \times n}_+$ if $A \otimes x = \lambda x$ for some $x \in \mathbb{R}^n_+ \setminus \{0\}$. The greatest eigenvalue of A will be denoted by $\lambda(A)$.

The vector $x \in \mathbb{R}^n_+ \setminus \{0\}$ satisfying $A \otimes x = \lambda x$ is called a *(max-algebraic)* eigenvector associated with A.

The eigencone of A associated with eigenvalue λ is defined as the set containing all eigenvectors of A with associated eigenvalue λ as well as the zero vector:

 $V(A,\lambda) := \{ x \in \mathbb{R}^n_+ \colon A \otimes x = \lambda \otimes x \}.$

One of the key notions of the paper is that of attraction cone: the set which comprises all vectors whose orbit hits a given eigencone.

Definition 2.2 (Attraction cones). The attraction cone of $A \in \mathbb{R}^{n \times n}_+$ associated with eigenvalue λ is the set

attr
$$(A, \lambda) = \{ x \in \mathbb{R}^n_+ : \mathcal{O}(A, x) \cap V(A, \lambda) \neq \emptyset \}.$$

We also denote $\operatorname{attr}(A) := \operatorname{attr}(A, \lambda(A)).$

Any eigencone or any attraction cone is a max cone, in the sense of the following definition.

Definition 2.3 (Max cones). A set $V \subseteq \mathbb{R}^n_+$ is called a *max cone* if for all $x \in V$, $y \in V$ any *max-linear combination* $\alpha x \oplus \beta y$ (where $\alpha, \beta \in \mathbb{R}_+$) belongs to V.

We will use the following notational shortcuts.

Definition 2.4 (Index Sets N and N_0). We denote

 $N = \{1, \dots, n\}, \quad N_0 = \{0, \dots, n-1\}.$

In this paper we deal with the following special class of matrices in maxalgebra.

Definition 2.5 (Circulant Matrices). A matrix $A \in \mathbb{R}^{n \times n}_+$ is called circulant, if it has entries $A_{i,j} = a_t$ for $i, j \in N, t \in N_0$ such that $t \equiv (j-i) \pmod{n}$ and $a_0, a_1, \ldots, a_{n-1} \in \mathbb{R}_+$. Equivalently, A is a circulant matrix if it is of the form

A =	$\int a_0$	a_1	a_2	 a_{n-2}	$a_{n-1} \\ a_{n-2}$	
	a_{n-}	$-1 a_0$	a_1	 a_{n-3}	a_{n-2}	
	:	÷	÷	÷	÷	
	a_1	a_2	a_3	 a_{n-1}	a_0)

for some $a_0, a_1, \ldots, a_{n-1} \in \mathbb{R}_+$. Such a circulant matrix we will also denote by $\mathcal{Z}(a_0, \ldots, a_{n-1})$.

Circulant matrices will be the main topic of Section 3 and Section 5, where we will study their spectral theory and attraction cones.

The final part of this paper is devoted to intervals and interval circulant matrices.

Definition 2.6 (Intervals). A set $X \subseteq \mathbb{R}^n_+$ is called an *interval* if it is of the form

$$\boldsymbol{X} = \times_{i=1}^{n} \boldsymbol{X}_{i},$$

for X_i nonempty subsets of \mathbb{R}_+ taking any of the following four forms:

$$[\underline{x}_i, \overline{x}_i], (\underline{x}_i, \overline{x}_i), (\underline{x}_i, \overline{x}_i], [\underline{x}_i, \overline{x}_i),$$

for $\underline{x}_i, \overline{x}_i \in \mathbb{R}_+$.

Definition 2.7 (Interval Circulant Matrices). By $\mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$ we denote the set of all circulant matrices A such that $A_{i,j} \in \boldsymbol{a}_{t}$ for $i, j \in N$ and $t \in N_{0}$ such that $t \equiv (j - i) \pmod{n}$, where $\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1}$ are intervals independently taking any of the four forms listed in Definition 2.6.

A set of circulant matrices that is of the form $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ for intervals $\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1}$ is called an interval circulant matrix.

In the literature on max-algebra, $A \in \mathbb{R}^{n \times n}_+$ is called robust if $x \in \operatorname{attr}(A)$ for all $x \in \mathbb{R}^n_+$, see [4] Section 8.6. In this paper we consider various extensions of this notion to interval circulant matrices. These extensions are listed in the following definition.

Definition 2.8 (Interval Robustness). Let $X \subseteq \mathbb{R}^n_+$ be an interval and $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ be an interval circulant matrix. Then $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ is called

possibly \mathbf{X} -robust if $(\exists A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))(\forall x \in \mathbf{X})[x \in \operatorname{attr}(A)]$, universally \mathbf{X} -robust if $(\forall A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))(\forall x \in \mathbf{X})[x \in \operatorname{attr}(A)]$, tolerance \mathbf{X} -robust if $(\forall A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))(\exists x \in \mathbf{X})[x \in \operatorname{attr}(A)]$, weakly tolerance \mathbf{X} -robust if $(\exists A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))(\exists x \in \mathbf{X})[x \in \operatorname{attr}(A)]$ and \mathbf{X} is called possibly \mathcal{Z}^{C} -robust if $(\exists x \in \mathbf{X})(\forall A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))[x \in \operatorname{attr}(A)]$,

tolerance \mathcal{Z}^C -robust if $(\forall x \in \mathbf{X})(\exists A \in \mathcal{Z}^C(\mathbf{a}_0, \dots, \mathbf{a}_{n-1}))[x \in \operatorname{attr}(A)]$.

In particular, the aim of Section 4 will be to derive an efficient characterization of these types of interval robustness.

2.2. Associated graphs, critical graphs and periodicity

Let us start with the following basic definition. For relevant definitions see also, e.g., [4] Section 1.5.

Definition 2.9 (Digraphs, Walks, Cycles and Connectivity). Let \mathcal{G} be a digraph with set of nodes N and set of edges E. A walk on \mathcal{G} is a sequence $W = (i_0, i_1, \ldots, i_l)$ with $i_0, i_1, \ldots, i_l \in N$ where each pair (i_{s-1}, i_s) for $s \in \{1, \ldots, l\}$ is an edge. If $i_0 = i$ and $i_l = j$ then W is said to be connecting i to j, and l is called the length of W.

 \mathcal{G} is called *strongly connected* if for each $i, j \in N$ with $i \neq j$ there exists a walk on \mathcal{G} connecting i to j.

For $A \in \mathbb{R}^{n \times n}_+$, the weighted digraph $\mathcal{G}(A)$ associated with A is the digraph with set of nodes $N = \{1, \ldots, n\}$ and set of edges $E = \{(i, j) : A_{i,j} \neq 0\}$, where $A_{i,j}$ is the weight of an edge (i, j).

If $\mathcal{G} = \mathcal{G}(A)$ then the weight of $W = (i_0, i_1, \dots, i_l)$ is defined by $A_{i_0, i_1} \cdot A_{i_1, i_2} \cdot \dots \cdot A_{i_{l-1}, i_l}$. This walk is called a cycle if $i_l = i_0$, with the cycle (geometric) mean defined by $(A_{i_0, i_1} \cdot A_{i_1, i_2} \cdot \dots \cdot A_{i_{l-1}, i_0})^{1/l}$.

Let us also give a separate definition of the maximum cycle mean.

Definition 2.10 (Maximum cycle (geometric) mean). The maximum cycle (geometric) mean of any $A \in \mathbb{R}^{n \times n}_+$ or of $\mathcal{G}(A)$ is

$$\max_{k=1}^{n} \max_{1 \le i_1, \dots, i_k \le n} (A_{i_1, i_2} \cdot A_{i_2, i_3} \dots A_{i_k, i_1})^{1/k}.$$
 (2)

The striking importance of this concept in max-algebra is due to the following fact.

Proposition 2.11 (e.g., [4], Corollary 4.5.6). For any $A \in \mathbb{R}^{n \times n}_+$, its greatest max-algebraic eigenvalue $(\lambda(A))$ is equal to (2).

The concept of irreducible matrix is common for max-algebra and nonnegative linear algebra, and it is most conveniently defined via the associated digraph.

Definition 2.12 (Irreducible, Reducible and Completely Reducible). A is called *irreducible* if $\mathcal{G}(A)$ is strongly connected, and *reducible* otherwise.

Digraph \mathcal{G} is called *completely reducible* if it consists of several strongly connected subgraphs called *components* such that there are no walks connecting a node from one component to a node of another component. A is called *completely reducible* if so is $\mathcal{G}(A)$.

Note that any irreducible matrix is completely reducible. Observe also the following criterion of complete reducibility.

Proposition 2.13. A digraph $\mathcal{G} = (N, E)$ is completely reducible if and only if every edge of E lies in a cycle of \mathcal{G} .

Proof. "If": Suppose that \mathcal{G} contains two maximal strongly connected subgraphs \mathcal{G}_1 and \mathcal{G}_2 and that there is a walk connecting one subgraph to the other. Without loss of generality we can assume that the walk does not contain nodes from any other subgraphs, so that it contains an edge (i, j) with $i \in \mathcal{G}_1$ and $j \in \mathcal{G}_2$. As this edge is on a cycle, there is also a walk from jto i. However, this implies that \mathcal{G}_1 and \mathcal{G}_2 both belong to a larger strongly connected subgraph of \mathcal{G} thus contradicting their maximality. Thus the "if" part is proved.

"Only if": If \mathcal{G} is completely reducible then each edge (i, j) belongs to a strongly connected subgraph of \mathcal{G} , and it belongs to a cycle since there exists a walk connecting j back to i.

The following subdigraph of $\mathcal{G}(A)$ is crucial for the max-algebraic spectral theory and it is an example of completely reducible digraph.

Definition 2.14 (Critical Digraphs). The critical digraph of A, denoted by $\mathcal{G}_c(A)$, consists of all nodes and edges of the cycles of $\mathcal{G}(A)$ at which the maximum cycle mean of A (2) is attained. These cycles are called critical cycles. The nodes of $\mathcal{G}_c(A)$ are called critical nodes and their set is denoted by $N_c(A)$, and the edges of $\mathcal{G}_c(A)$ are called critical edges and their set is denoted by $E_c(A)$.

Corollary 2.15. Any critical graph is completely reducible.

Proof. By Definition 2.14, every edge of $\mathcal{G}_c(A)$ belongs to a cycle of $\mathcal{G}_c(A)$. The claim now follows from Proposition 2.13.

The concept of the digraph's cyclicity is crucial for the study of attraction cones (Definition 2.2) and the ultimate periodicity of $\{A^t\}_{t\geq 1}$ (to be defined soon).

Definition 2.16 (Cyclicity). For a strongly connected digraph, its *cyclic-ity* is defined as the g.c.d. of the lengths of all cycles of that digraph.

Cyclicity of a completely reducible digraph is defined as the l.c.m. of the cyclicities of its components.

Cyclicity of a digraph \mathcal{G} is denoted by $\sigma(\mathcal{G})$.

We now discuss the ultimate periodicity of max-algebraic matrix powers.

Definition 2.17 (Ultimate Periodicity). Let $\{\alpha_k\}_{k\geq 1}$ be a sequence of some elements. If there exists T such that $\alpha_{t+\sigma} = \alpha_t$ for all $t \geq T$ and some σ (i.e., $\alpha_{t+\sigma}$ and α_t are identical), then $\{\alpha_k\}_{k\geq 1}$ is called *ultimately periodic*. The least T and the least σ for which the above property holds are called the *transient* and the *ultimate period* of $\{\alpha_k\}_{k>1}$ respectively.

Proposition 2.18 ([5]). Let $A \in \mathbb{R}^{n \times n}_+$ be an irreducible matrix with $\lambda(A) \neq 0$. Then $\{(A/\lambda(A))^t\}_{t\geq 1}$ is ultimately periodic and $\sigma(\mathcal{G}_c(A))$ is the ultimate period of that sequence.

In this paper we also need the following trivial extension of Proposition 2.18 and its consequence for orbits of vectors.

Corollary 2.19. Let $A \in \mathbb{R}^{n \times n}_+$ be a completely reducible matrix with $\lambda(A) \neq 0$, such that the maximum cycle mean of each component of $\mathcal{G}(A)$ is the same (and equal to $\lambda(A)$). Then $\{(A/\lambda(A))^t\}_{t\geq 1}$ is ultimately periodic and $\sigma(\mathcal{G}_c(A))$ is the ultimate period of that sequence.

Corollary 2.20. Under the conditions of Proposition 2.18 or Corollary 2.19, $\{(A/\lambda(A))^t \otimes x\}_{t \ge 1}$ is ultimately periodic for any $x \in \mathbb{R}^n_+$.

Let us now introduce some notation related to the ultimate periodicity.

Definition 2.21. Let $A \in \mathbb{R}^{n \times n}_+$ have $\lambda(A) \neq 0$. If $\{(A/\lambda(A))^t\}_{t \geq 1}$ is ultimately periodic then denote by T(A) the transient and by per(A) the ultimate period of that sequence.

Thus $per(A) = \sigma(\mathcal{G}_c(A))$ for any A satisfying the condition of Proposition 2.18 or Corollary 2.19.

The ultimate period of $\{(A/\lambda(A))^t \otimes x\}_{t \ge 1}$ does not necessarily equal the cyclicity of $\mathcal{G}_c(A)$, and the attraction cone associated with $\lambda(A)$ consists of the vectors for which the ultimate period of $\{(A/\lambda(A))^t \otimes x\}_{t \ge 1}$ is equal to 1. More precisely, we have the following.

Proposition 2.22. Let $A \in \mathbb{R}^{n \times n}_+$ be a completely reducible matrix with $\lambda(A) \neq 0$ such that the maximum cycle mean of each component of $\mathcal{G}(A)$ is the same (and equal to $\lambda(A)$). Then

$$\operatorname{attr}(A) = \{ x \in \mathbb{R}^n_+ \colon \lambda(A) A^t \otimes x = A^{t+1} \otimes x \}, \quad where \ t \ge T(A).$$

Proof. By definition $x \in \operatorname{attr}(A)$ if and only if $A^{s+1} \otimes x = \lambda(A)A^s \otimes x$, hence $\lambda(A)A^t \otimes x = A^{t+1} \otimes x$ is sufficient for $x \in \operatorname{attr}(A)$. For the necessity observe that $A^{s+1} \otimes x = \lambda(A)A^s \otimes x$ implies $A^{s'+1} \otimes x = \lambda(A)A^{s'} \otimes x$ for some $s' \geq \max(s, T(A))$ and such that $(A/\lambda(A))^{s'} = (A/\lambda(A))^t$, and hence $A^{t+1} \otimes x = \lambda(A)A^t \otimes x$.

Corollary 2.23. Under the conditions of Proposition 2.22 $\operatorname{attr}(A)$ is a closed max-cone.

Proof. Under these conditions $\operatorname{attr}(A)$ is the solution set of the system $\lambda(A)A^t \otimes x = A^{t+1} \otimes x$. This solution set is a max-cone since it is closed under taking max-linear combinations (see Definition 2.3) and it is a closed set since all arithmetic operations of max-algebra are continuous.

Let us finally consider the attraction cones of the following two matrices satisfying the conditions of Proposition 2.22.

Example 2.24. Take

$$A = \begin{pmatrix} 0.5 & 1 & 0.2 & 0 \\ 1 & 0.5 & 0.2 & 0 \\ 0.2 & 0.2 & 0.2 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, B = \begin{pmatrix} 0.5 & 1 & 0.2 & 0 \\ 1 & 0.5 & 0.3 & 0 \\ 0.4 & 0.4 & 0.4 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

The ultimate periods of $\{A, A^2, A^3, \ldots\}$ and $\{B, B^2, B^3, \ldots\}$ equal 2. In the first case, the periodicity starts from A^2 (i.e., we have $A^2 = A^4$), and in the second case it starts from B^3 (i.e., we have $B^3 = B^5$). The attraction cones are

$$\operatorname{attr}(A) = \{ x \colon A^3 \otimes x = A^4 \otimes x \}, \quad \operatorname{attr}(B) = \{ x \colon B^3 \otimes x = B^4 \otimes x \},$$

where

$$A^{3} = \begin{pmatrix} 0.5 & 1 & 0.2 & 0 \\ 1 & 0.5 & 0.2 & 0 \\ 0.2 & 0.2 & 0.04 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \quad A^{4} = \begin{pmatrix} 1 & 0.5 & 0.2 & 0 \\ 0.5 & 1 & 0.2 & 0 \\ 0.2 & 0.2 & 0.04 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix},$$
$$B^{3} = \begin{pmatrix} 0.5 & 1 & 0.2 & 0 \\ 1 & 0.5 & 0.3 & 0 \\ 0.4 & 0.4 & 0.12 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \quad B^{4} = \begin{pmatrix} 1 & 0.5 & 0.3 & 0 \\ 0.5 & 1 & 0.2 & 0 \\ 0.4 & 0.4 & 0.12 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}.$$

We further see that in both cases, the systems defining these attraction cones reduce to just one equation:

attr(A) = {x: 0.5x₁
$$\oplus$$
 x₂ \oplus 0.2x₃ = x₁ \oplus 0.5x₂ \oplus 0.2x₃},
attr(B) = {x: 0.5x₁ \oplus x₂ \oplus 0.2x₃ = x₁ \oplus 0.5x₂ \oplus 0.3x₃},

Observe that $x = \begin{bmatrix} 1 & 1 & 5 & 1 \end{bmatrix}$ belongs to $\operatorname{attr}(A)$ but not to $\operatorname{attr}(B)$, and $x = \begin{bmatrix} 0.5 & 1 & \frac{10}{3} & 1 \end{bmatrix}$ belongs to $\operatorname{attr}(B)$ but not to $\operatorname{attr}(A)$.

Example 2.24 also shows that Theorem 3.10, the main result of the next section which claims that $\operatorname{attr}(A) \subseteq \operatorname{attr}(B)$ for any circulant A, B with $A \leq B$ and $\lambda(A) = \lambda(B)$, is not true for general completely reducible (or irreducible) matrices.

3. Circulant matrices: critical graph and attraction cones

Let us start with the following statement, which is well known in usual linear algebra. See, e.g., [6] Theorem 3.1.1. A proof of it in max-algebra, which works equally well in the usual linear algebra case, is given below for the reader's convenience.

Proposition 3.1. Let $A, B \in \mathbb{R}^{n \times n}_+$ be circulant matrices. Then $A \otimes B$ is also circulant. In particular, any max-algebraic power of A (or B) is a circulant.

Proof. Observe that A is a circulant matrix if and only if we can represent $A = a_0 I \oplus a_1 P \oplus \ldots \oplus a_{n-1} P^{n-1}$, where

$$P = \begin{pmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \dots & \ddots & \dots & \vdots \\ 0 & 0 & \dots & 0 & 1 \\ 1 & 0 & \dots & 0 & 0 \end{pmatrix}.$$

In this case $A = \mathcal{Z}(a_0, a_1, \ldots, a_{n-1})$. Computing $A \otimes B$ amounts to multiplying $a_0I \oplus a_1P \oplus \ldots \oplus a_{n-1}P^{n-1}$ by $b_0I \oplus b_1P \oplus \ldots \oplus b_{n-1}P^{n-1}$, assuming

that $A = \mathcal{Z}(a_0, a_1, \ldots, a_{n-1})$ and $B = \mathcal{Z}(b_0, b_1, \ldots, b_{n-1})$. This multiplication results in an expression of the form $c_0 I \oplus c_1 P \oplus \ldots \oplus c_{n-1} P^{n-1}$, thus also a circulant.

Writing A^t as $A^{t-1} \otimes A$ for every $t \ge 2$, we also show that A^t is a circulant by a simple inductive argument.

The following observation will play a key role in proving many properties of circulants.

Lemma 3.2. Let $A \in \mathbb{R}^{n \times n}_+$ be a nonzero circulant matrix and let $A_{i,j} = \mu \neq 0$ for some $i, j \in N$. Then (i, j) belongs to a cycle (i_1, \ldots, i_n, i_1) with $i_t - i_{t-1} \equiv (j-i) \pmod{n}$ for all $t \in \{2, \ldots, n\}$, and $i_1 - i_n \equiv (j-i) \pmod{n}$. The weight of each edge in (i_1, \ldots, i_n, i_1) equals $A_{i,j} = \mu$.

Proof. Consider an infinite sequence $\{i_\ell\}_{\ell\geq 1}$ where $i_1 = i$, $i_2 = j$, $i_{\ell+1} - i_\ell \equiv (j-i) \pmod{n}$ for all $\ell \geq 1$ and $i_\ell \in N$ for all $\ell \geq 1$. By Definition 2.5, $A_{i_\ell,i_{\ell+1}} = A_{i,j}$ for all $\ell \geq 1$. However, we also have that $i_{n+1} = i_1$ since $i_{n+1} - i_1 \equiv n \cdot (j-i) \pmod{n} = 0 \pmod{n}$. Hence the claim follows. \Box

Proposition 3.3. Let $A = \mathcal{Z}(a_0, \ldots, a_{n-1})$. Then A has a unique maxalgebraic eigenvalue equal to

$$\lambda(A) = \max_{k=0}^{n-1} a_k. \tag{3}$$

If $A \neq 0$ then $\lambda(A) \neq 0$ and all nodes in N are critical.

Proof. If $A \neq 0$ then $\max(a_0, \ldots, a_{n-1}) > 0$. In this case, let *i* and *j* be such that $A_{i,j} = \mu > 0$. By Lemma 3.2 (i, j) belongs to a cycle (i_1, \ldots, i_n, i_1) where the weights of all edges are equal to μ . It follows that the cycle mean of that cycle is also μ . Thus, the maximal cycle mean is equal to the maximal weight of edges, which shows (3). Taking *k* such that $a_k = \lambda(A)$, for each $i \in N$ we have *j* with $k \equiv (j - i) \pmod{n}$ such that $A_{i,j} = \lambda(A)$, hence each $i \in N$ is on a critical cycle. Since all nodes $\mathcal{G}(A)$ are critical, *A* has a unique eigenvalue equal to $\lambda(A)$ as it follows, e.g., from [4] Corollary 4.5.8.

If A = 0 then $\max(a_0, \dots, a_{n-1}) = 0 = \lambda(A)$.

Note that equation (3) was obtained already in [16], Theorem 2.1. However, we preferred to give a partially self-contained proof of this equation for the reader's convenience. **Corollary 3.4.** Let $A \in \mathbb{R}^{n \times n}_+$ be a circulant matrix. Then $\lambda(A) = 0$ if and only if A = 0.

Proof. Obviously, $\lambda(A) = 0$ if A = 0. The "only if" part is equivalent to the implication $(A \neq 0) \Rightarrow (\lambda(A) \neq 0)$ stated in Proposition 3.3.

We now formulate the following immediate corollary of Proposition 2.13.

Corollary 3.5. Any circulant matrix A is completely reducible.

Proof. If A = 0 then $\mathcal{G}(A)$ has no edges and is completely reducible. Otherwise, by Lemma 3.2 any edge of $\mathcal{G}(A)$ belongs to a cycle, and the claim follows from Proposition 2.13.

Proposition 3.6. For any nonzero circulant matrix $A \in \mathbb{R}^{n \times n}_+$ the matrix sequence $\{(A/\lambda(A))^t\}_{t\geq 1}$ is ultimately periodic, and $T(A) \leq (n-1)^2 + 1$.

Proof. For the first part of the claim observe that any circulant matrix is completely reducible by Corollary 3.5, and that by Proposition 3.3 $\lambda(A)$ is the maximum cycle mean of any maximal strongly connected component of $\mathcal{G}(A)$.

Since $\lambda(A) = 0$ implies A = 0 by Corollary 3.4, we can assume $\lambda(A) = 1$ without loss of generality. Since all nodes of $\mathcal{G}(A)$ are critical, the transient of periodicity of $\{A, A^2, A^3, \ldots\}$ is the same as the greatest transient of periodicity of any sequence of rows of these powers $\{A_{i\bullet}, A_{i\bullet}^2, A_{i\bullet}^3, \ldots\}$ where *i* is critical. However, these transients are bounded by $(n-1)^2 + 1$ by [13] Main Theorem 1.

The following proposition gives more information on the critical graph and cyclicity of circulant matrices.

Proposition 3.7. Let $A = \mathcal{Z}(a_0, \ldots, a_{n-1}) \neq 0$ and let $p_1, \ldots, p_s \in \{1, \ldots, n-1\}$ be the nonzero indices for which $a_{p_1} = \ldots = a_{p_s} = \lambda(A)$ (if such indices exist) and such that $p_1 > p_2 > \ldots > p_s$. Then

(i) $\mathcal{G}_c(A)$ consists of $m = \gcd(n, p_1, \ldots, p_s)$ isomorphic strongly connected components. Node set of the *i*th component, for $i \in \{1, \ldots, m\}$, is $\{i, i + m, \ldots, i + (n/m - 1)m\}$.

(ii) per(A), equal to the cyclicity of each of these components, is 1 if $a_0 = \lambda(A)$ and

$$per(A) = gcd(\frac{n}{gcd(n, p_1)}, \frac{p_1 - p_2}{gcd(p_1, p_2)}, \frac{p_1 - p_3}{gcd(p_1, p_3)}, \dots, \frac{p_1 - p_s}{gcd(p_1, p_s)})$$

$$= gcd(\frac{n}{gcd(n, p_1)}, \frac{p_1 - p_2}{gcd(p_1, p_2)}, \frac{p_2 - p_3}{gcd(p_2, p_3)}, \dots, \frac{p_{s-1} - p_s}{gcd(p_{s-1}, p_s)})$$

$$= gcd(\frac{n}{gcd(n, p_1)}, \frac{p_1 - p_2}{gcd(n, p_1, p_2)}, \frac{p_1 - p_3}{gcd(n, p_1, p_2, p_3)}, \dots, \frac{p_1 - p_s}{gcd(n, p_1, p_2, p_3)})$$

$$(4)$$

if
$$a_0 \neq \lambda(A)$$
.

Parts of this statement can be found in [15] Theorem 4.1 and Lemma 4.1. Essentially, part (i) was proved in [14] Lemma 4.2 and Lemma 4.3, although in the max-min algebra setting. The number $gcd(n, p_1, \ldots, p_s)$ also appeared in [22] Theorem 4 as the "eigenspace dimension". The result of part (ii) relies on [9] (Theorems 3.1 and 3.3) where the cyclicity of threshold circulant graphs (see Definition 5.3) was studied. We will give a complete proof of (i) and a reduction of (ii) to the results of [9] in Subsection 5.1, for the reader's convenience.

Let us now describe the attraction cone of a circulant matrix as a solution set of a max-algebraic two-sided system of equations.

Proposition 3.8. Let $A \in \mathbb{R}^{n \times n}_+$ be a circulant matrix. Then

$$\operatorname{attr}(A) = \{ x \colon \lambda(A)A^{n^2} \otimes x = A^{n^2+1} \otimes x \}.$$

Proof. By Corollary 3.4, $\lambda(A) = 0$ if and only if A = 0, in which case $\operatorname{attr}(A) = \mathbb{R}^n_+$ and $A^{n^2} = A^{n^2+1} = 0$, and the claim holds trivially. Otherwise, by Corollary 3.5 A is completely reducible and by Proposition 3.3 the maximal cycle mean of each component of $\mathcal{G}(A)$ is the same. The claim then follows since A satisfies the conditions of Proposition 2.22 and since $n^2 \geq T(A)$ by Proposition 3.6.

Let us examine the attraction cone of a 4×4 circulant matrix.

Example 3.9. Consider

$$A = \begin{pmatrix} 0 & 0 & 1 & t \\ t & 0 & 0 & 1 \\ 1 & t & 0 & 0 \\ 0 & 1 & t & 0 \end{pmatrix},$$

where t: 0 < t < 1. This is a circulant matrix, $\lambda(A) = 1$, and $\mathcal{G}_c(A)$ consists of two disjoint cycles: (1 3) and (2 4). The cyclicity of $\mathcal{G}_c(A)$ is thus equal to 2 and so is the ultimate period of the max-algebraic matrix powers of A. Taking the max-algebraic powers of A we obtain

$$A^{2} = \begin{pmatrix} 1 & t & t^{2} & 0 \\ 0 & 1 & t & t^{2} \\ t^{2} & 0 & 1 & t \\ t & t^{2} & 0 & 1 \end{pmatrix}, \quad A^{2k} = \begin{pmatrix} 1 & t & t^{2} & t^{3} \\ t^{3} & 1 & t & t^{2} \\ t^{2} & t^{3} & 1 & t \\ t & t^{2} & t^{3} & 1 \end{pmatrix} \quad \forall k \ge 2,$$
$$A^{2k-1} = \begin{pmatrix} 0 & 0 & 1 & t \\ t & 0 & 0 & 1 \\ 1 & t & 0 & 0 \\ 0 & 1 & t & 0 \end{pmatrix} \quad \forall k \ge 2.$$

In particular, the periodicity transient is T(A) = 3. By Proposition 3.8 we have $\operatorname{attr}(A) = \{x \colon A^{16} \otimes x = A^{17} \otimes x\}$, implying that the attraction cone is precisely the set of vectors $x = (x_1 \ x_2 \ x_3 \ x_4)$ that satisfy

$$x_1 \oplus tx_2 \oplus t^2 x_3 \oplus t^3 x_4 = t^2 x_1 \oplus t^3 x_2 \oplus x_3 \oplus tx_4$$

$$tx_1 \oplus t^2 x_2 \oplus t^3 x_3 \oplus x_4 = t^3 x_1 \oplus x_2 \oplus tx_3 \oplus t^2 x_4.$$
(5)

System (5) can be further reduced using the cancellation rule

$$a \oplus b = ta \oplus c \Leftrightarrow a \oplus b = c,$$

where t < 1 and a, b, c are arbitrary. Repeatedly applying this rule we obtain the system

$$\begin{aligned} x_1 \oplus tx_2 &= x_3 \oplus tx_4 \\ tx_1 \oplus x_4 &= x_2 \oplus tx_3, \end{aligned} \tag{6}$$

equivalent to (5).

Now observe that $x = [t \ 1 \ t^2 \ 1]$ satisfies this system of equations and belongs to the attraction cone. In particular, the ultimate period of $\{A^t x\}_{t\geq 1}$ is 1, however, $A \otimes x \neq x$ which shows that $\operatorname{attr}(A)$ is not the same as the (max-algebraic) eigencone of A in this case. The following theorem is one of the main results of the paper. Its proof is postponed to Subsection 5.2.

Theorem 3.10. Let $A, B \in \mathbb{R}^{n \times n}_+$ be two circulant matrices such that $\lambda(A) = \lambda(B)$ and $A \leq B$. Then $\operatorname{attr}(A) \subseteq \operatorname{attr}(B)$.

Let us give two examples demonstrating this theorem. In the first example we have two 0-1 matrices, and in the second one we consider the matrix of Example 3.9 with two different values of t.

Example 3.11. Let us first consider a pair of 0-1 matrices:

	(0)	1	0	0	0	0			$\begin{pmatrix} 0 \\ 0 \end{pmatrix}$	1	0	1	0	0	
A =	$\begin{pmatrix} 0 \\ 0 \end{pmatrix}$	0	1	0	0	0	,	, B =							
	0	0	0	1	0	0			0	0	0	1	0	1	
	0	0	0	0	1	0			1	0	0	0	1	0	
	0	0	0	0	0	1			0						
	$\begin{pmatrix} 0\\ 1 \end{pmatrix}$	0	0	0	0	0/			$\backslash 1$	0	1	0	0	0/	

Observe that the sequence $\{A^t\}_{t\geq 1}$ is periodic from the very beginning. The system $A^{36} \otimes x = A^{37} \otimes x$, being the same as $A \otimes x = A^2 \otimes x$, reduces to $x_1 = x_2 = x_3 = x_4 = x_5 = x_6$.

The sequence $\{B^t\}_{t\geq 1}$ becomes periodic from T(B) = 3. More precisely, we have

$$B^{2k-1} = \begin{pmatrix} 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \end{pmatrix}, \quad B^{2k} = \begin{pmatrix} 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}, \quad k \ge 2,$$

and the system $B^{36} \otimes x = B^{37} \otimes x$ reduces to $x_1 \oplus x_3 \oplus x_5 = x_2 \oplus x_4 \oplus x_6$, thus $\operatorname{attr}(A) \subseteq \operatorname{attr}(B)$.

Example 3.12. Take

$$A = \begin{pmatrix} 0 & 0 & 1 & t_1 \\ t_1 & 0 & 0 & 1 \\ 1 & t_1 & 0 & 0 \\ 0 & 1 & t_1 & 0 \end{pmatrix}, \quad B = \begin{pmatrix} 0 & 0 & 1 & t_2 \\ t_2 & 0 & 0 & 1 \\ 1 & t_2 & 0 & 0 \\ 0 & 1 & t_2 & 0 \end{pmatrix}$$

with $0 < t_1 < t_2 < 1$. Then $\operatorname{attr}(A)$ is the set of all x satisfying (6) with $t = t_1$, which is

$$\begin{aligned} x_1 \oplus t_1 x_2 &= x_3 \oplus t_1 x_4 \\ t_1 x_1 \oplus x_4 &= x_2 \oplus t_1 x_3, \end{aligned} \tag{7}$$

and $\operatorname{attr}(B)$ is the set of all x satisfying

$$x_1 \oplus t_2 x_2 = x_3 \oplus t_2 x_4$$

$$t_2 x_1 \oplus x_4 = x_2 \oplus t_2 x_3,$$
(8)

We next show that $\operatorname{attr}(A) \subseteq \operatorname{attr}(B)$ in this example, by considering various special cases.

Suppose first that we have $t_1x_2 = t_1x_4 \ge x_1 \oplus x_3$ in the first equation of (7). This implies $x_2 = x_4 \ge t_2(x_1 \oplus x_3) \ge t_1(x_1 \oplus x_3)$ and $t_2x_2 = t_2x_4 \ge$ $(x_1 \oplus x_3)$. This shows that in this case x belongs to both attr(A, 1) and attr(B, 1). The case when $t_1x_1 = t_1x_3 \ge x_2 \oplus x_4$ in the second equation of (7) is treated similarly.

Suppose now that $x \in \operatorname{attr}(A)$ and $t_1x_2 = x_3 \ge x_1 \oplus t_1x_4$. As we cannot have $t_1x_1 = x_2$ and $x_4 = t_1x_3$ in the second equation of (7), assume that $x_2 = x_4 \ge t_1(x_1 \oplus x_3)$. But this implies $t_1x_2 = t_1x_4$, and as t_1x_2 is the maximum in the first equation, this returns us to the case which we considered first, where $x \in \operatorname{attr}(B)$. We also note three other similar cases that are treated in the same way.

The remaining case when $x \in \operatorname{attr}(A)$, $x_1 = x_3 \ge t_1(x_2 \oplus x_4)$ and $x_2 = x_4 \ge t_1(x_1 \oplus x_3)$ is impossible when $t_1 < 1$.

4. Interval robustness of circulant matrices

In this section we characterize the six types of interval robustness of Definition 2.8 for interval circulant matrix $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ and interval $\boldsymbol{X} = \times_{i=1}^{n} \boldsymbol{X}_{i}$ where \boldsymbol{X}_{i} and \boldsymbol{a}_{i} are intervals independently taking one of the following four forms:

$$[\underline{x}_i, \overline{x}_i], (\underline{x}_i, \overline{x}_i), (\underline{x}_i, \overline{x}_i], [\underline{x}_i, \overline{x}_i)$$

and

$$[\underline{a}_j, \overline{a}_j], \ (\underline{a}_j, \overline{a}_j), \ (\underline{a}_j, \overline{a}_j], \ [\underline{a}_j, \overline{a}_j)$$

for $\underline{x}_i, \overline{x}_i \in \mathbb{R}_+$ and $i \in N$, and $\underline{a}_j, \overline{a}_j \in \mathbb{R}_+$ and $j \in N_0$, respectively.

4.1. Universal and possible X-robustness

Let us introduce the following notation.

Definition 4.1 (Matrices $A^{(k)}$ and vectors $x^{(k)}$). For a given index $k \in N_0$ denote

$$A^{(k)} = \mathcal{Z}(\underline{a}_0, \underline{a}_1, \dots, \underline{a}_{k-1}, \overline{a}_k, \underline{a}_{k+1}, \dots, \underline{a}_{n-1}),$$

and

$$x^{(k)} = (\underline{x}_1, \underline{x}_2, \dots, \underline{x}_{k-1}, \overline{x}_k, \underline{x}_{k+1}, \dots, \underline{x}_n)$$

The following lemma explains the use of vectors $x^{(k)}$.

Lemma 4.2. Let $\mathbf{X} \subseteq \mathbb{R}^n_+$ be an interval and let $A \in \mathbb{R}^{n \times n}_+$. Then $\mathbf{X} \subseteq \operatorname{attr}(A)$ if and only if $x^{(i)} \in \operatorname{attr}(A)$ for each $i \in N$.

Proof. Observe first that since the cone $\operatorname{attr}(A)$ is a closed set by Corollary 2.23, the inclusion $\mathbf{X} \subseteq \operatorname{attr}(A)$ is equivalent to $\operatorname{cl}(\mathbf{X}) \subseteq \operatorname{attr}(A)$, where cl is a Euclidean closure. Since $x^{(i)} \in \operatorname{cl}(\mathbf{X})$ for all $i \in N$ (as vertices of the box $\operatorname{cl}(\mathbf{X})$), it follows that the condition is necessary. Let us show that this condition is also sufficient. For this we will show that

$$x = \bigoplus_{k=1}^{n} \frac{x_k}{\overline{x}_k} x^{(k)}.$$
(9)

Indeed, observe that when $k \neq i$ we have that $x_k/\overline{x}_k \leq 1$ implies $(x_k/\overline{x}_k)x_i^{(k)} \leq \underline{x}_i$, and when k = i we obtain $(x_i/\overline{x}_i)x_i^{(i)} = x_i$. Since $\underline{x}_i \leq x_i$, we obtain that

$$\bigoplus_{k=1}^{n} \frac{x_k}{\overline{x}_k} x_i^{(k)} = (x_i/\overline{x}_i) x_i^{(i)} = x_i,$$

for all i, so (9) holds. Thus x can be expressed as a max-linear combination of $x^{(k)}$ for $k \in N$ and $x \in \operatorname{attr}(A)$ since $\operatorname{attr}(A)$ is a max-cone (Definition 2.3). \Box

Definition 4.3 (Matrix \hat{A}). For $\underline{a} = \max_{k \in N_0} \underline{a}_k$ define

$$\hat{A} = \mathcal{Z}(\hat{a}_0, \hat{a}_1, \dots, \hat{a}_{n-1}),$$

where

$$\hat{a}_i = \min\{\underline{a}, \overline{a}_i\}, \text{ for each } i \in N_0.$$

Let us characterize the cases when $\hat{A} = 0$ and when $\hat{A} \in \mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$.

Proposition 4.4. Let $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ be given. Then

- (i) $\hat{A} = 0 \Leftrightarrow \underline{a} = 0 \Leftrightarrow \underline{A} = 0 \Leftrightarrow \lambda(\underline{A}) = 0.$
- (ii) If $\forall i : \mathbf{a}_i = [\underline{a_i}, \overline{a_i}]$, then $\hat{A} \in \mathcal{Z}^C(\mathbf{a}_0, \dots, \mathbf{a}_{n-1})$.
- (iii) $\hat{A} \in \mathcal{Z}^C(\boldsymbol{a}_0, \dots, \boldsymbol{a}_{n-1}) \Leftrightarrow \forall i : (\underline{a} \ge \overline{a}_i \Rightarrow \overline{a}_i \in \boldsymbol{a}_i) \& (\underline{a} \le \overline{a}_i \Rightarrow \underline{a} \in \boldsymbol{a}_i).$

Proof. (i): Let us show that $\hat{A} = 0 \Leftrightarrow \underline{a} = 0$. By Definition 4.3 it is immediate that $\underline{a} = 0$ implies $\hat{A} = 0$. Next, assume that $\hat{A} = 0$. Then $\overline{a}_i = 0$ for all i, which implies $\underline{a}_i = 0$ for all i, hence $\underline{a} = 0$. The equivalence $\underline{a} = 0 \Leftrightarrow \underline{A} = 0$ is obvious, and $\underline{A} = 0 \Leftrightarrow \lambda(\underline{A}) = 0$ follows from Corollary 3.4. (ii) and (iii): Straightforward.

Matrices \hat{A} and $A^{(k)}$ for k = 0, ..., n - 1 have the following useful properties.

Lemma 4.5. If
$$\hat{A} \neq 0$$
, then $(\forall A \in \mathcal{Z}^C(\boldsymbol{a}_0, \dots, \boldsymbol{a}_{n-1}))[(A/\lambda(A)) \leq (\hat{A}/\lambda(\hat{A})]$

Proof. Observe that $\hat{A} \neq 0$ implies that A = 0 does not belong to the interval matrix $\mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$. Recalling that $\hat{a}_{i} = \min(\overline{a}_{i}, \underline{a})$ for all i we see that $\hat{a}_{i} \leq \underline{a}$ for all i and that $\hat{a}_{k} = \underline{a}$ for k such that $\underline{a}_{k} = \underline{a}$. Hence $\lambda(\hat{A}) = \underline{a}$ by Proposition 3.3. Showing $(A/\lambda(A)) \leq (\hat{A}/\lambda(\hat{A}))$ means showing

$$a_i / \max_k a_k \le \min(\overline{a}_i, \underline{a}) / \underline{a} \quad \forall i.$$
 (10)

To prove (10) we observe that it follows from the inequality

$$a_i \cdot \underline{a} \le \max_j a_j \cdot \min(\overline{a}_i, \underline{a}) \quad \forall i,$$
(11)

which is

$$a_i \cdot \underline{a} \le \max_i a_j \cdot \underline{a} \tag{12}$$

when $\min(\overline{a}_i, \underline{a}) = \underline{a}$, and

$$a_i \cdot \max_i \underline{a}_i \le \overline{a}_i \cdot \max_j a_j \tag{13}$$

when $\min(\overline{a}_i, \underline{a}) = \overline{a}_i$. Both (12) and (13) are obvious. This shows (11) and hence (10) and $(A/\lambda(A)) \leq (\hat{A}/\lambda(\hat{A}))$.

Lemma 4.6. For any nonzero $A \in \mathbb{Z}^C(\mathbf{a}_0, \ldots, \mathbf{a}_{n-1})$ there exists $A^{(k)} \neq 0$ for some $k \in N_0$ such that $[(A^{(k)}/\lambda(A^{(k)})) \leq (A/\lambda(A))].$

Proof. Let $A = \mathcal{Z}(a_0, \ldots, a_{n-1})$ and let k be such that $a_k = \max_{j \in N} a_j$. Consider $A^{(k)}$. Since $\overline{a}_k \ge a_k > 0$ but the rest of the components defining $A^{(k)}$ are $\underline{a}_i \le a_i$ for $i \ne k$, we have $\lambda(A^{(k)}) = \overline{a}_k$ and $(A^{(k)}/\lambda(A^{(k)})) \le (A/\lambda(A))$. \Box

We now characterize possibly X-robust and universally X-robust interval circulant matrices.

Theorem 4.7. Let $X \subseteq \mathbb{R}^n_+$ be an interval, and let $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix containing \hat{A} . Then $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ is possibly X-robust if and only if we have $x^{(i)} \in \operatorname{attr}(\hat{A})$ for all $i \in N$.

Proof. We need to show that there exists $A \in \mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$ such that $\boldsymbol{X} \subseteq \operatorname{attr}(A)$ if and only if $x^{(i)} \in \operatorname{attr}(\hat{A})$ for all $i \in N$. If $\hat{A} = 0$ then $\operatorname{attr}(\hat{A}) = \mathbb{R}^{n}_{+}$ and the claim is obvious. Next we suppose that $\hat{A} \neq 0$ which implies $\lambda(\hat{A}) \neq 0$ by Corollary 3.4. By Proposition 4.4 part (i), $\mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$ contains only nonzero matrices in this case.

"If": By Lemma 4.2, the condition implies that $X \subseteq \operatorname{attr}(\hat{A})$. The claim then follows since $\hat{A} \in \mathbb{Z}^{\mathbb{C}}(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$.

"Only if": Let $A \in \mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$ be such that $\boldsymbol{X} \subseteq \operatorname{attr}(A)$. By Lemma 4.5 we have $(A/\lambda(A)) \leq (\hat{A}/\lambda(\hat{A}))$, and Theorem 3.10 yields that $x \in \operatorname{attr} \hat{A}$. As $x \in \operatorname{attr} \hat{A}$ for all $x \in \boldsymbol{X}$, the claim then follows from Lemma 4.2.

Corollary 4.8. Let $x \in \mathbb{R}^n_+$ and let $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix containing \hat{A} . Then $(\exists A \in \mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}))[x \in \operatorname{attr}(A)]$ if and only if $x \in \operatorname{attr}(\hat{A})$.

Proof. Take $X = \{x\}$ then the possible X-robustness means existence of $A \in \mathbb{Z}^{\mathbb{C}}(a_0, \ldots, a_{n-1})$ such that $x \in \operatorname{attr}(A)$ and $x^{(i)} = x$ for all $i \in N$. The claim then follows from Theorem 4.7.

Theorem 4.9. Let $X \subseteq \mathbb{R}^n_+$ be an interval, and let $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix. Then $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ is universally X-robust if and only if $x^{(j)} \in \operatorname{attr}(A^{(i)})$ for all $i \in N_0$ and $j \in N$.

Proof. We need to show that $X \subseteq \operatorname{attr}(A)$ for all $A \in \mathbb{Z}^{C}(a_{0}, \ldots, a_{n-1})$ if and only if $x^{(j)} \in \operatorname{attr}(A^{(i)})$ for all $i \in N_{0}$ and $j \in N$.

"If": Let $x^{(j)} \in \operatorname{attr}(A^{(i)})$ hold for all $i \in N_0$ and $j \in N$. Take $A \in \mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$. If A = 0 then $x^{(j)} \in \operatorname{attr}(A) = \mathbb{R}^n_+$. Otherwise, by Lemma 4.6 there exists $k \in N_0$ such that $A^{(k)} \neq 0$ and $(A^{(k)}/\lambda(A^{(k)})) \leq (A/\lambda(A))$. Applying Theorem 3.10 to $(A^{(k)}/\lambda(A^{(k)}))$ and $(A/\lambda(A))$ we obtain $x^{(j)} \in \operatorname{attr}(A)$ for all nonzero $x^{(j)}$, hence $\mathbf{X} \subseteq \operatorname{attr}(A)$.

"Only if": Take a sequence $\{A_s\}_{s\geq 1} \subseteq \mathcal{Z}^C(\boldsymbol{a}_0,\ldots,\boldsymbol{a}_{n-1})$ such that $\lim_{s\to\infty} A_s = A^{(k)}$, and take any $x \in \boldsymbol{X}$. Since $x \in \operatorname{attr}(A_s)$ for all s, by Proposition 3.8 we have $\lambda(A_s)A_s^{n^2} \otimes x = A_s^{n^2+1} \otimes x$ for all s, and by the continuity of the arithmetic operations of max-algebra we obtain $\lambda(A^{(k)})(A^{(k)})^{n^2} \otimes x = (A^{(k)})^{n^2+1} \otimes x$. As $x \in \operatorname{attr}(A^{(k)})$ for all $x \in \boldsymbol{X}$, the claim then follows from Lemma 4.2.

Corollary 4.10. Let $x \in \mathbb{R}^n_+$, and let $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix. Then $(\forall A \in \mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})) [x \in \operatorname{attr}(A)]$ if and only if $x \in \operatorname{attr}(A^{(k)})$ for each $k \in N_0$.

Proof. Take $\mathbf{X} = \{x\}$ then the universal \mathbf{X} -robustness means that $x \in \operatorname{attr}(A)$ for all $A \in \mathbb{Z}^{C}(\mathbf{a}_{0}, \ldots, \mathbf{a}_{n-1})$. The claim then follows from Theorem 4.9.

4.2. Tolerance and weak tolerance X-robustness

Theorem 4.11. Let $X \subseteq \mathbb{R}^n_+$ be a closed interval, and let $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix. Then $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ is tolerance X-robust if and only if $(\forall k \in N_0)[(\operatorname{attr}(A^{(k)}) \cap X) \neq \emptyset].$

Proof. "If": Take $A \in \mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$. If A = 0 then $\operatorname{attr}(A) = \mathbb{R}^n_+$, hence $\operatorname{attr}(A) \cap \boldsymbol{X} \neq \emptyset$. Otherwise, for each $i \in N_0$ take $y^{(i)} \in (\boldsymbol{X} \cap \operatorname{attr}(A^{(i)})]$, By Lemma 4.6 there exists $k \in N_0$ with $(A^{(k)}/\lambda(A^{(k)})) \leq (A/\lambda(A))$. Applying Theorem 3.10 to $(A^{(k)}/\lambda(A^{(k)}))$ and $(A/\lambda(A))$ we obtain $y^{(k)} \in \operatorname{attr}(A)$, hence the implication.

"Only if": For any $k \in N_0$ take a sequence $\{A_s\}_{s\geq 1} \subseteq \mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ such that $\lim_{s\to\infty} A_s = A^{(k)}$. For each of these matrices there exists $x^s \in \boldsymbol{X}$ such that $x^s \in \operatorname{attr}(A_s)$. Then by Proposition 3.8 we have $\lambda(A_s)A_s^{n^2} \otimes x^s = A_s^{n^2+1} \otimes x^s$ for all s. Since \boldsymbol{X} is compact, we can assume that $\lim_{s\to\infty} x^s$ exists

and denote it by $y^{(k)}$. Then we obtain that by the continuity of arithmetic operations of max-algebra $\lambda(A^{(k)})(A^{(k)})^{n^2} \otimes y^{(k)} = (A^{(k)})^{n^2+1} \otimes y^{(k)}$. Hence $y^{(k)} \in \operatorname{attr}(A^{(k)})$.

Corollary 4.12. Under the conditions of Theorem 4.11, $\mathcal{Z}^{C}(\boldsymbol{a}_{0}, \ldots, \boldsymbol{a}_{n-1})$ is tolerance \boldsymbol{X} -robust if and only if all systems

$$\lambda(A^{(k)})(A^{(k)})^{n^2} \otimes y = (A^{(k)})^{n^2+1} \otimes y, \quad y \in \boldsymbol{X},$$
(14)

with $k \in N_0$ such that $A^{(k)} \neq 0$ are solvable.

We now characterize the weak tolerance robust matrices.

Theorem 4.13. Let $\mathbf{X} \subseteq \mathbb{R}^n_+$ be an interval and let $\mathcal{Z}^C(\mathbf{a}_0, \ldots, \mathbf{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix containing \hat{A} . Then $\mathcal{Z}^C(\mathbf{a}_0, \ldots, \mathbf{a}_{n-1})$ is weakly tolerance \mathbf{X} -robust if and only if $\lambda(\hat{A})(\hat{A})^{n^2} \otimes x = (\hat{A})^{n^2+1} \otimes x$ is solvable with $x \in \mathbf{X}$.

Proof. By Corollary 4.8, $x \in \mathbf{X}$ and $A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \ldots, \mathbf{a}_{n-1})$ such that $x \in \operatorname{attr}(A)$ exist if and only if $x \in \operatorname{attr}(\hat{A})$ for some $x \in \mathbf{X}$. This, by Proposition 3.8, is equivalent to $\lambda(\hat{A})(\hat{A})^{n^{2}} \otimes x = (\hat{A})^{n^{2}+1} \otimes x$ being solvable with $x \in \mathbf{X}$.

4.3. Possible and tolerance $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ -robustness

We now characterize the remaining two types of robustness.

Theorem 4.14. Let $\mathbf{X} \subseteq \mathbb{R}^n_+$ be an interval, and let $\mathcal{Z}^C(\mathbf{a}_0, \ldots, \mathbf{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix. Then \mathbf{X} is possibly $\mathcal{Z}^C(\mathbf{a}_0, \ldots, \mathbf{a}_{n-1})$ -robust if and only if there exists $x \in \mathbf{X}$ that satisfies $\lambda(A^{(i)})(A^{(i)})^{n^2} \otimes x = (A^{(i)})^{n^2+1} \otimes x$ for all $i \in N_0$ such that $A^{(i)} \neq 0$.

Proof. By Corollary 4.10, $x \in \mathbf{X}$ belongs to attr(A) for all $A \in \mathcal{Z}^{C}(\mathbf{a}_{0} \dots, \mathbf{a}_{n-1})$ if and only if it belongs to attr($A^{(i)}$) for all $i \in N_{0}$ with $A^{(i)} \neq 0$. By Proposition 3.8 this is equivalent to x satisfying $\lambda(A^{(i)})(A^{(i)})^{n^{2}} \otimes x = (A^{(i)})^{n^{2}+1} \otimes x$ for all such i. **Theorem 4.15.** Let $X \subseteq \mathbb{R}^n_+$ be an interval, and let $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1}) \subseteq \mathbb{R}^{n \times n}_+$ be an interval circulant matrix containing \hat{A} . Then interval vector X is tolerance $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ -robust if and only if $\mathcal{Z}^C(\boldsymbol{a}_0, \ldots, \boldsymbol{a}_{n-1})$ is possibly X-robust.

Proof. Suppose that X is tolerance $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ -robust, then we have the following

$$(\forall x \in \mathbf{X})(\exists A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))[x \in \operatorname{attr}(A)] \stackrel{\operatorname{Cor.4.8}}{\Longleftrightarrow} (\forall x \in \mathbf{X})[x \in \operatorname{attr}(\hat{A})]$$
$$\Rightarrow (\exists A \in \mathcal{Z}^{C}(\mathbf{a}_{0}, \dots, \mathbf{a}_{n-1}))(\forall x \in \mathbf{X})[x \in \operatorname{attr}(A)],$$

and hence we have that $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ is possibly \boldsymbol{X} -robust.

The converse implication is trivial.

4.4. Computational complexity

We close the section with a couple of remarks on the computational complexity of the different types of interval robustness.

Remark 4.16. By Theorems 4.7 and 4.15 the verification of whether

- (i) $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ is possibly **X**-robust,
- (ii) $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ is universally \boldsymbol{X} -robust,
- (iii) \boldsymbol{X} is tolerance $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ -robust

reduces, under some assumptions, to the verification whether some vectors satisfy some two-sided max-linear systems with n^2 and $n^2 + 1$ powers of some matrices. Hence these types of robustness are of polynomial complexity.

Remark 4.17. By Corollary 4.12, Theorem 4.14 and Theorem 4.15, verifying whether

- (i) $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ is tolerance \boldsymbol{X} -robust,
- (ii) $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ is weakly tolerance \boldsymbol{X} -robust,
- (iii) \boldsymbol{X} is possibly $\mathcal{Z}^{C}(\boldsymbol{a}_{0},\ldots,\boldsymbol{a}_{n-1})$ -robust

reduces, under some assumptions, to verifying the non-emptyness of solution set of some system of max-affine inequalities, where some of the inequalities (among those defining \mathbf{X}) can be strict. This problem was generally shown to be polynomially equivalent to solving a mean-payoff game [1], for which efficient pseudopolynomial algorithms exist, but existence of a polynomial algorithm has been a long-standing open question.

5. Proofs of Proposition 3.7 and Theorem 3.10

5.1. Cyclicity of circulants: Proof of Proposition 3.7

Let us start with the following elementary but useful statement.

Lemma 5.1. Let $p_1, \ldots, p_s, n \in \mathbb{N}$ (the set of natural numbers). Then the equation

$$p_1 x_1 + \ldots + p_s x_s \equiv m \pmod{n} \tag{15}$$

has a solution $x = (x_1, \ldots, x_s) \in \mathbb{N}^s$ if and only if m is a multiple of $gcd(p_1, \ldots, p_s, n)$.

Proof. "Only if": Observe that $p_1x_1 + \ldots + p_sx_s$ and n are always multiples of $gcd(p_1, \ldots, p_s, n)$, and if (15) holds then so is m as well.

"If": The claim is well known for s = 1 (elementary number theory). The same fact also implies existence of $x_s \in \mathbb{N}$ such that

$$p_s x_s \equiv m \pmod{\gcd(n, p_1, \dots, p_{s-1})}.$$
(16)

We now prove the claim by induction assuming that it holds for s-1. Observe that (16) implies that there exists also $k \in \mathbb{N}$ such that

$$p_s x_s + k \gcd(n, p_1, \dots, p_{s-1}) \equiv m \pmod{n}. \tag{17}$$

But by induction there exist $x_1 \in \mathbb{N}, \ldots, x_{s-1} \in \mathbb{N}$ such that

$$p_1x_1 + \ldots + p_{s-1}x_{s-1} \equiv k \gcd(n, p_1, \ldots, p_{s-1}) \pmod{n}.$$
 (18)

Combining (17) and (18) we get the claim.

Let us now introduce the following definition that appeared in [19] (see also [4]).

Definition 5.2 (Visualized Matrices). A nonzero $A \in \mathbb{R}^{n \times n}_+$ is called

- (i) visualized if $A_{i,j} \leq \lambda(A)$ for all i, j, and
- (ii) strictly visualized if it is visualized and $A_{i,j} = \lambda(A)$ if and only if $(i, j) \in \mathcal{G}_c(A)$.

By (3) we have that $\lambda(A) = \max(a_0, a_1, \dots, a_{n-1})$ for $A = \mathcal{Z}(a_0, \dots, a_{n-1})$, implying that $\lambda(A) = \max_{i,j=1}^n A_{i,j}$ for any circulant A. That is, any circulant matrix is visualized. We will now argue that it is also strictly visualized.

Definition 5.3 (Threshold Digraphs). Let $A \in \mathbb{R}^{n \times n}_+$ and $h \in \mathbb{R}_+$. Define the *threshold digraph* of A with respect to h as the subgraph of $\mathcal{G}(A)$ containing all edges (i, j) with $A_{i,j} \geq h$, and all nodes that are beginning and end nodes of those edges. Denote this threshold graph by $\mathcal{G}(A, h)$.

Proposition 5.4. Let $A \in \mathbb{R}^{n \times n}_+$ be a nonzero circulant matrix. Then it is strictly visualised, and $\mathcal{G}_c(A) = \mathcal{G}(A, \lambda(A))$.

Proof: By (3) no entry of A exceeds $\lambda(A)$. Hence A is visualized. Also recall that $\lambda(A) > 0$ by Corollary 3.4.

If $A_{i,j} < \lambda(A)$ then the mean weight of any cycle with edge (i, j) is strictly less than $\lambda(A)$, so (i, j) is not critical. In other words, (i, j) being critical implies $A_{i,j} = \lambda(A)$.

It remains to show that if $A_{i,j} = \lambda(A)$, which is equivalent to (i, j) being an edge of $\mathcal{G}(A, \lambda(A))$, then (i, j) is critical. In this case by Lemma 3.2 (i, j)lies in a cycle with all edge weights equal to $\lambda(A)$. The weights of all edges in this cycle are equal to $\lambda(A)$, hence the mean weight of this cycle is $\lambda(A)$, i.e., it is a critical cycle and (i, j) is critical. This completes the proof. \Box

Proof of Proposition 3.7. First observe that Proposition 5.4 implies that $\mathcal{G}_c(A) = \mathcal{G}(A, \lambda(A))$ and hence the set of critical edges of a circulant matrix A is given by

$$E_{c}(A) = \{(i, j) : i = j \text{ if } a_{0} = \lambda(A) \text{ or } j - i \equiv p_{k} (\text{mod } n), k \in \{1, \dots, s\}\}$$
(19)
where p_{1}, \dots, p_{s} are such that $a_{p_{1}} = \dots = a_{p_{s}} = \lambda(A) (\text{and } p_{1} > p_{2} > \dots > p_{s}).$

We now consider the component of $\mathcal{G}_c(A)$ which contains node *i*, for *i* from the set $\{1, \ldots, \gcd(n, p_1, \ldots, p_s)\}$.

Let us argue that the node set of this component is given by

$$\{k \in N : k \equiv i + l_1 p_1 + \ldots + l_s p_s (\text{mod } n), \ l_1, \ldots, l_s \in \mathbb{N} \cup \{0\}\},$$
(20)

Indeed, by (19) edges (i, j) where $j \equiv (l+p_t) \pmod{n}$ for some $t \in \{1, \ldots, s\}$ are the only edges that issue from i and are critical. Using this observation, the claim follows by simple induction.

Using Lemma 5.1 we now observe that (20) is the same as

$$\{i + k \operatorname{gcd}(n, p_1, \dots, p_s)\}: k \in \{0, \dots, (n/\operatorname{gcd}(n, p_1, \dots, p_s)) - 1\}.$$

This set does not intersect with the node set of any component containing a different node in $\{1, \ldots, \gcd(n, p_1, \ldots, p_s)\}$, and this yields $\gcd(n, p_1, \ldots, p_s)$ strongly connected components of $\mathcal{G}_c(A)$. Isomorphism between two components containing $i_1 \in \{1, \ldots, \gcd(n, p_1, \ldots, p_s)\}$ and $i_2 \in \{1, \ldots, \gcd(n, p_1, \ldots, p_s)\}$ is induced by the following mapping on their set of nodes:

$$i_1 + k \operatorname{gcd}(n, p_1, \dots, p_s)) \mapsto i_2 + k \operatorname{gcd}(n, p_1, \dots, p_s)).$$

This completes the proof of part (i) of Proposition 3.7.

If $a_0 = \lambda(A)$ then $\mathcal{G}_c(A)$ contains all loops of the form (i, i) for $1 \leq i \leq n$, and the cyclicity of every component of $\mathcal{G}_c(A)$ is 1 since it contains a loop. When $a_0 < \lambda(A)$, we can use the result of [9] Theorem 3.3 part (i) since this result describes the cyclicity of any component of the threshold digraph $\mathcal{G}(A, \lambda(A))$ (see [9] Theorem 3.1.), and since $\mathcal{G}_c(A) = \mathcal{G}(A, \lambda(A))$ by Proposition 5.4. According to this result, that cyclicity is equal to any of the three expressions given in (4). This completes the proof of part (ii).

5.2. Inclusion of attraction cones: Proof of Theorem 3.10

Before considering the problem of our interest, let us recall the notion of cyclic classes which will be necessary for some proofs.

Definition 5.5 (Cyclic Classes). Let $\mathcal{G} = (N, E)$ be a strongly connected graph with cyclicity $\sigma(\mathcal{G})$, and let $i, j \in N$. Nodes i, j are said to belong to the same *cyclic class* if the lengths of some (and hence all) walks connecting i to j are a multiple of $\sigma(\mathcal{G})$.

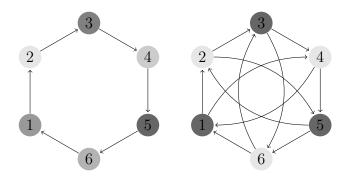


Figure 1: Cyclic classes of two graphs of Example 5.6 (shown in different shades).

The cyclic class of i will be denoted by [i]. We also write $[i] \rightarrow_1 [j]$ if the lengths of some (and hence all) walks connecting a member of [i] to a member of [j] have length congruent to 1 modulo $\sigma(\mathcal{G})$.

By cyclic classes of a completely reducible digraph we mean cyclic classes of its (strongly connected) components.

Example 5.6. Consider two associated graphs of 0-1 matrices of Example 3.11 shown in Figure 1. On the left, the graph consists just of one cycle of length 6, hence its cyclicity is 6 and the cyclic classes are $\{1\}$, $\{2\}$, $\{3\}$, $\{4\}$, $\{5\}$ and $\{6\}$. On the right, the cyclicity of the graph is 2 and the cyclic classes are $\{1, 3, 5\}$ and $\{2, 4, 6\}$.

Cyclic classes are also called *components of imprimitivity*. We refer the reader to [3] Lemma 3.4.1 for a proof that belonging to the same cyclic class is a well-defined equivalence relation.

Lemma 5.7. Let \mathcal{G} be a strongly connected digraph.

- (i) Let $\sigma(\mathcal{G}) > 1$ and let i_0, i_1, \ldots, i_k be a walk on \mathcal{G} . Then $[i_{l-1}] \rightarrow_1 [i_l]$ for each $l \in \{1, \ldots, k\}$.
- (ii) Let C be a cycle of G. Then C contains a member of each cyclic class of G.

Proof. (i): Each edge is a walk of length 1. Therefore $[i_{l-1}] \rightarrow_1 [i_l]$ for each $l \in \{1, \ldots, k\}$.

(ii): Let i be a node which is not in C. Let us show that C contains a node

in the cyclic class of i. Since \mathcal{G} is strongly connected, there exists a walk connecting i to a node j of C. If the length of this walk is a multiple of $\sigma(\mathcal{G})$ then $j \in [i]$. Otherwise, we concatenate this walk with a walk from j to some node $k \in C$ whose edges belong to C and such that the length of resulting walk is a multiple of $\sigma(\mathcal{G})$. Then $k \in [i]$ and the claim is proved. \Box

We now derive a convenient form of a system defining the attraction cone for circulant matrices, based on the results of [18]. Here $A_{i\bullet}^t$ denotes the *i*th row of A^t . We also write $i \sim_A j$ when *i* and *j* belong to the same component of $\mathcal{G}_c(A)$.

Proposition 5.8. Let $A \in \mathbb{R}^{n \times n}_+$ be a nonzero circulant matrix. Then

$$x \in \operatorname{attr}(A) \Leftrightarrow A_{i\bullet}^{n^2} \otimes x = A_{j\bullet}^{n^2} \otimes x \; \forall i, j \in N \, s.t. \, [i] \to_1 [j]$$
(21)

and

$$x \in \operatorname{attr}(A) \Leftrightarrow A_{i\bullet}^{n^2} \otimes x = A_{j\bullet}^{n^2} \otimes x \; \forall i, j \in N \text{ s.t. } i \sim_A j.$$

$$(22)$$

Proof. By Proposition 3.8

$$\operatorname{attr}(A) = \{ x \colon A^{n^2} \otimes x = A^{n^2 + 1} \otimes x \}$$
(23)

Since A is a circulant matrix, by Proposition 5.4 it is visualized, and then by [18] Proposition 2.8 we also have $A_{j\bullet}^{n^2} = A_{i\bullet}^{n^2+1}$ for any $i, j \in N_c(A)$ such that $[i] \rightarrow_1 [j]$. This shows (21). To show (22) recall that if a component of $\mathcal{G}_c(A)$ has more than one cyclic class then for every two nodes i, j of the component there is a walk $i_0 = i, i_1, i_2, \ldots, i_k = j$ on $\mathcal{G}_c(A)$ where $[i_{l-1}] \rightarrow_1 [i_l]$ for each $l \in \{1, \ldots, k\}$ by Lemma 5.7 part (i). Hence $A_{i\bullet}^{n^2} \otimes x = A_{j\bullet}^{n^2} \otimes x$ holds for all nodes i, j in that component. If a component has only one cyclic class then [18] Proposition 2.8 implies that all rows with indices in that component are equal to each other, so the equations $A_{i\bullet}^{n^2} \otimes x = A_{j\bullet}^{n^2} \otimes x$ hold trivially for all pairs of nodes from that component. \Box

It can be seen that we wrote out system (21) for all examples of Section 3. In the case of Example 3.11, for which $\mathcal{G}(A) = \mathcal{G}_c(A)$ and the cyclic classes are shown on Figure 1, system (21) reduces to $x_1 = x_2 = x_3 = x_4 = x_5 = x_6$ for A and to $x_1 \oplus x_3 \oplus x_5 = x_2 \oplus x_4 \oplus x_6$ for B.

We will also need the following observations.

Lemma 5.9. Let $A, B \in \mathbb{R}^{n \times n}_+$ be two matrices such that $\lambda(A) = \lambda(B) \neq 0$ and $A \leq B$. Then $\mathcal{G}_c(A) \subseteq \mathcal{G}_c(B)$.

Proof. Since $A \leq B$ the mean weight of each cycle in B is not less than the mean weight of the same cycle in A. If that cycle is critical in A then its mean weight $\lambda(A)$ cannot increase in B since $\lambda(A) = \lambda(B)$. Hence it equals $\lambda(B)$, i.e., the cycle belongs to $\mathcal{G}_c(B)$.

Lemma 5.10. Let $A, B \in \mathbb{R}^{n \times n}_+$ be two circulant matrices with $\lambda(A) = \lambda(B) \neq 0, A \leq B$. Then

$$x \in \operatorname{attr}(B) \Leftrightarrow B_{i\bullet}^{n^2} \otimes x = B_{j\bullet}^{n^2} \otimes x \quad \forall i, j \in N, s.t. \ i \sim_A j.$$
(24)

Proof. By (22),

$$x \in \operatorname{attr}(B) \Leftrightarrow B_{i\bullet}^{n^2} \otimes x = B_{j\bullet}^{n^2} \otimes x \ \forall i, j \in N \text{ s.t. } i \sim_B j.$$
 (25)

We also have $\mathcal{G}_c(A) \subseteq \mathcal{G}_c(B)$ by Lemma 5.9 and hence each $x \in \operatorname{attr}(B)$ satisfies the system in (24).

Suppose now that x satisfies the system in(24). We will show that x also satisfies

$$B_{i\bullet}^{n^2} \otimes x = B_{j\bullet}^{n^2} \otimes x \; \forall i, j \in N \text{ s.t. } [i] \to_1 [j]$$

$$(26)$$

so that $x \in \operatorname{attr}(B)$ by Proposition 5.8. Since $\mathcal{G}_c(A) \subseteq \mathcal{G}_c(B)$, each component α of $\mathcal{G}_c(A)$ belongs to a component β of $\mathcal{G}_c(B)$, and each component of $\mathcal{G}_c(B)$ contains a component of $\mathcal{G}_c(A)$ because $N_c(A) = N_c(B) = N$. Hence it amounts to show that if x satisfies the subsystem of equations in (24) corresponding to a component α of $\mathcal{G}_c(A)$ then it also satisfies the subsystem of equations in (26) corresponding to the component β of $\mathcal{G}_c(B)$ such that $\alpha \subseteq \beta$. But by Lemma 5.7 part (ii) each cyclic class of β has a member in any cycle of β and hence in any cycle of α (because $\alpha \subseteq \beta$). This shows that for each i, j with $[i] \to_1 [j]$ in $\mathcal{G}_c(B)$ there exist $k \in [i]$ and $l \in [j]$ on a cycle of α and then $B_{k\bullet}^{n^2} \otimes x = B_{l\bullet}^{n^2} \otimes x$ holds by (24). However, $B_{k\bullet}^{n^2} = B_{i\bullet}^{n^2}$ and $B_{l\bullet}^{n^2} = B_{j\bullet}^{n^2}$ by [18] Proposition 2.8. Hence the claim follows.

Let us now introduce Kleene stars, as they will also be useful in the proof of Theorem 3.10.

Definition 5.11 (Kleene Stars). Let $A \in \mathbb{R}^{n \times n}_+$ have $\lambda(A) \leq 1$. Then

$$A^* = I \oplus A \oplus A^2 \oplus \ldots \oplus A^{n-1}$$

is called the *Kleene star of A*.

Proposition 5.12 (e.g., [4], Corollary 1.6.16). Let $A \in \mathbb{R}^{n \times n}_+$. Then $A^* = A$ if and only if one of the following equivalent conditions hold:

- (i) $A^2 = A$ and $A_{i,i} = 1$ for all $i \in N$;
- (ii) $A_{i,i} = 1$ and $A_{i,j}A_{j,k} \leq A_{i,k}$ for all $i, j, k \in N$.

More specifically, we will make use of the following.

Lemma 5.13. Let $A \neq 0$ be a circulant matrix. Then $(A/\lambda(A))^{n^2}$ is a Kleene star.

Proof. Note that $\lambda(A) \neq 0$ by Corollary 3.4. By Proposition 5.12 it suffices to show that $(A/\lambda(A))^{n^2}$ is an idempotent matrix and that $((A/\lambda(A))^{n^2})_{i,i} = 1$ for all *i*. For the idempotency, observe that by Proposition 3.7 part (ii) per(A) divides n^2 , and that $T(A) \leq n^2$ by Proposition 3.6. Hence $(A/\lambda(A))^{2n^2} = (A/\lambda(A))^{n^2}$.

For the remaining part of the claim, assume $\lambda(A) = 1$ and recall that for any $t \geq 1$ and any $i, j \in N$, entry $(A^t)_{i,j}$ is equal to the greatest weight of a walk of length t connecting i to j (e.g.,[4], Example 1.2.3). Take $i \in \{1, \ldots, n\}$ and observe that $\mathcal{G}(A)$ contains a critical cycle of length ngoing through i. The weights of all entries of that cycle equal to 1. Taking n copies of this cycle we obtain a cycle in $\mathcal{G}(A)$ of weight 1 and length n^2 . The claim $((A/\lambda(A))^{n^2})_{i,i} = 1$ follows since the weights of all entries and (therefore) of all walks are bounded by 1.

We are now ready to prove the main result of Section 3.

The proof will also make use of the following notation.

Definition 5.14. Denote by $k \mod n$, respectively by $k \mod n$, the only number in $N_0 = \{0, \ldots, n-1\}$, respectively in $N = \{1, \ldots, n\}$, which is congruent to k modulo n.

Proof of Theorem 3.10. The case $\lambda(A) = \lambda(B) = 0$ is trivial since in that case A = B = 0 by Corollary 3.4 and hence $\operatorname{attr}(A) = \operatorname{attr}(B) = \mathbb{R}^n_+$. Otherwise, as $\operatorname{attr}(A/\lambda(A)) = \operatorname{attr}(A)$ and $\operatorname{attr}(B/\lambda(B)) = \operatorname{attr}(B)$ (which follows, e.g., from Proposition 2.22), we can assume without loss of generality $\lambda(A) = \lambda(B) = 1$ and consider matrices $C = A^{n^2}$ and $D = B^{n^2}$. By Proposition 3.1 C and D are circulants, hence $C = \mathcal{Z}(c_0, \ldots, c_{n-1})$ and D = $\mathcal{Z}(d_0, \ldots, d_{n-1})$ for some c_0, \ldots, c_{n-1} and d_0, \ldots, d_{n-1} . Using $A \oplus B = B$ and the expansion for $(A \oplus B)^{n^2}$ we obtain $A^{n^2} \leq (A \oplus B)^{n^2} = B^{n^2}$, thus $C \leq D$. By Lemma 5.13 both of them are also Kleene stars. By Proposition 5.12 we have $D_{1,(\alpha+\gamma)[\operatorname{mod}' n]} \geq D_{1,\alpha} \cdot D_{\alpha,(\alpha+\gamma)[\operatorname{mod}' n]}$ and hence

$$d_{(\alpha+\gamma-1)[\text{mod }n]} \ge d_{\alpha-1} \cdot d_{\gamma} \tag{27}$$

for any $\alpha \in \{1, \ldots, n\}$ and $\gamma \in \{0, \ldots, n-1\}$. In what follows we are going to show that the assumption that $\operatorname{attr}(A) \subseteq \operatorname{attr}(B)$ does not hold leads to a contradiction with (27) for some α and γ .

By Lemma 5.9 we have $\mathcal{G}_c(A) \subseteq \mathcal{G}_c(B)$. By Proposition 3.7 part (i), $\mathcal{G}_c(A)$ consists of l components whose node sets are of the form

$$\{k, k+l, k+2l, \dots, k+(n/l-1)l\} \text{ for } k \in \{1, \dots, l\},$$
(28)

where l is a divisor of n. Each of these node sets belongs to some component of $\mathcal{G}_c(B)$.

By Proposition 5.8 $x \in \operatorname{attr}(A)$ if and only if

$$C_{k\bullet} \otimes x = C_{k+l\bullet} \otimes x = \dots = C_{k+(n/l-1)l\bullet} \otimes x \quad \text{for } k \in \{1, \dots, l\},$$
(29)

and by Lemma 5.10 $x \in \operatorname{attr}(B)$ if and only if

$$D_{k\bullet} \otimes x = D_{k+l\bullet} \otimes x = \dots = D_{k+(n/l-1)l\bullet} \otimes x \quad \text{for } k \in \{1, \dots, l\}, \quad (30)$$

We will refer to (29) or (30) for fixed k as to a *chain of equations*.

Suppose by contradiction that $x \in \operatorname{attr}(A)$ but $x \notin \operatorname{attr}(B)$. The latter means that there exist k and s such that $D_{k\bullet} \otimes x > D_{k+ls\bullet} \otimes x$ for some integers k and s. Assume without loss of generality that k = 1 then

$$D_{1\bullet} \otimes x = d_0 x_1 \oplus d_1 x_2 \oplus \ldots \oplus d_{n-1} x_n.$$

Let $d_{\alpha-1} \cdot x_{\alpha}$ be one of the terms where the maximum in the above expression is attained. In $D_{1+ls\bullet} \otimes x$ we find a term $d_{\alpha-1} \cdot x_{\beta}$ where $\alpha \equiv \beta \pmod{l}$, and we have the inequality $d_{\alpha-1} \cdot x_{\alpha} > d_{\alpha-1} \cdot x_{\beta}$ and hence $x_{\alpha} > x_{\beta}$.

Observe that $c_0 = d_0 = 1$ since C and D are Kleene stars. Since $\alpha \equiv \beta \pmod{l}$ there exists a chain of equations among those of (29), which contains both $c_0 x_{\alpha} = x_{\alpha}$ and $c_0 x_{\beta} = x_{\beta}$. The corresponding chain of equations holds (since $x \in \operatorname{attr}(A)$), but $x_{\alpha} > x_{\beta}$ and therefore in the expression containing $c_0 x_{\beta}$ there is a term $c_{\gamma} x_{(\beta+\gamma)[\operatorname{mod}' n]}$ (for some γ) such that $c_{\gamma} x_{(\beta+\gamma)[\operatorname{mod}' n]} \ge x_{\alpha} > 0$, and hence

$$d_{\gamma} x_{(\beta+\gamma)[\text{mod}' n]} \ge x_{\alpha}. \tag{31}$$

Going back to the terms in the inequality $D_{1,\bullet}x > D_{1+ls,\bullet}x$ and knowing that the maximum in $D_{1,\bullet}x$ is attained at $d_{\alpha-1}x_{\alpha}$ and $D_{1+ls,\bullet}x$ contains a term of the form $d_{\alpha-1}x_{\beta}$, we see that $D_{1+ls,\bullet}x$ also contains the term $d_{(\alpha+\gamma-1)[\text{mod }n]}x_{(\beta+\gamma)[\text{mod' }n]}$ and that

$$d_{\alpha-1}x_{\alpha} > d_{(\alpha+\gamma-1)[\text{mod }n]}x_{(\beta+\gamma)[\text{mod' }n]}.$$
(32)

Multiplying (31) by $d_{\alpha-1}$, combining with (32) and canceling $x_{(\beta+\gamma)[\text{mod}' n]} > 0$ we have

$$d_{\alpha-1}d_{\gamma} > d_{(\alpha+\gamma-1)[\text{mod }n]},$$

which contradicts with the Kleene star property (27). The proof is complete. $\hfill \Box$

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